Time Warp: Global Control

Distributed Snapshots and Fossil Collection

Richard M. Fujimoto
Professor

Computational Science and Engineering Division
College of Computing
Georgia Institute of Technology
Atlanta, GA 30332-0765, USA

http://www.cc.gatech.edu/~fujimoto/

Outline

- Consistent Cuts
  - Cut points
  - Cut messages
  - Cut values
- Mattern's GVT Algorithm
  - Colors
  - Vector counters
  - Pipelined algorithm
- Fossil Collection

Mattern's Algorithm

- Asynchronous: executes in background concurrent with time warp execution (does not require the simulation to block)
- Avoids message acknowledgements
- Based on techniques for creating distributed snapshots

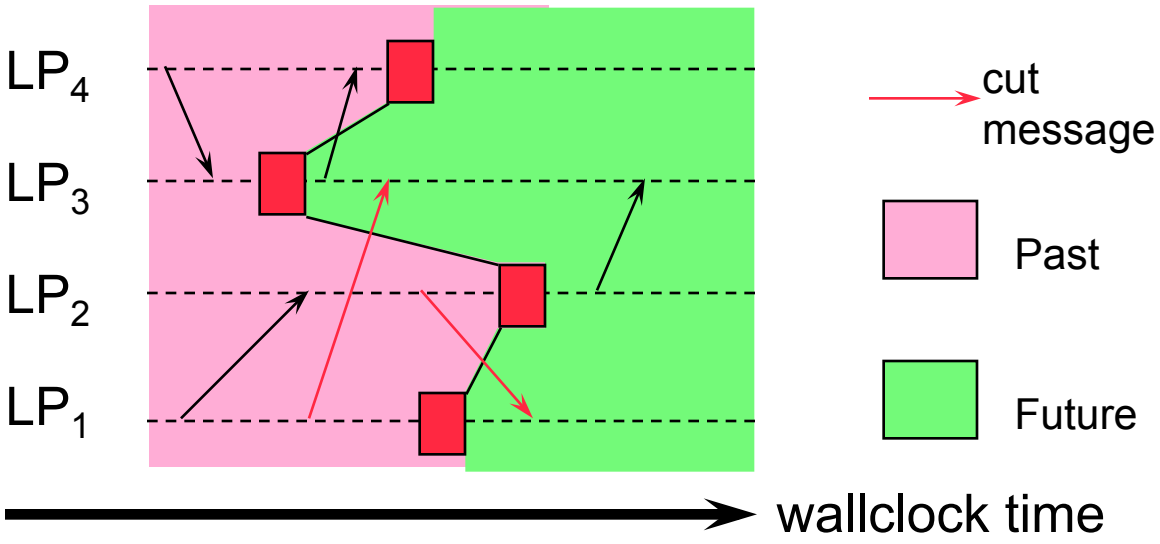
Mattern: Consistent Cuts

cut point: an instant dividing computation into past and future

cut: set of cut points, one per processor

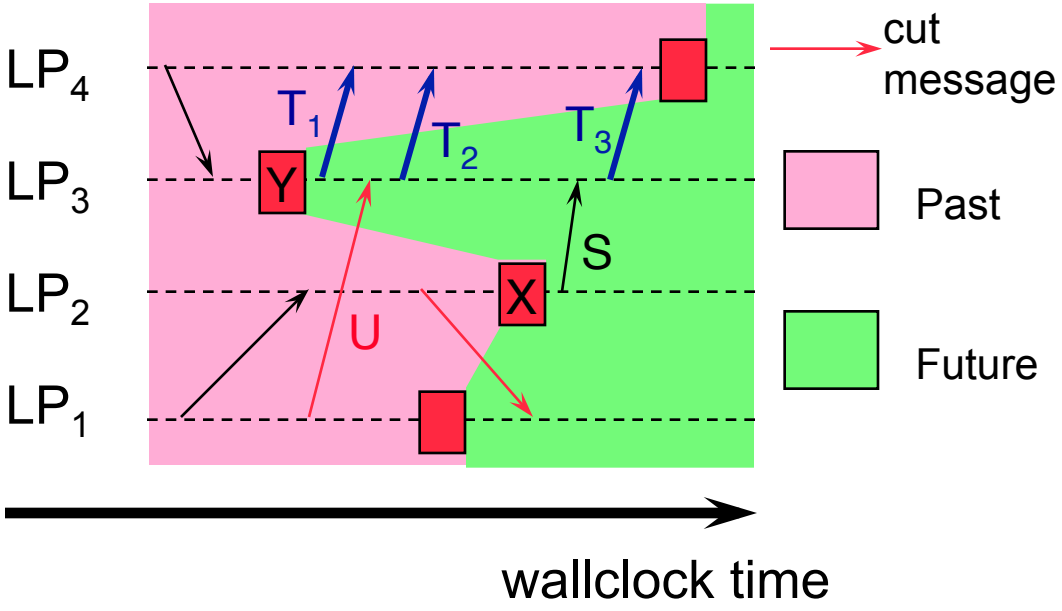
cut message: a message that was sent in the past, and received in the future

consistent cut: a cut where all messages crossing the cut are cut messages



cut value: minimum among (1) local minimum of each LP at its cut point and (2) time stamp of cut messages

Observation 1



$T_1 > Y$
$T_2 > \min(Y, U)$
$T_3 > \min(Y, U, S)$
$S > X$, so
$T_3 > \min(Y, U, X)$

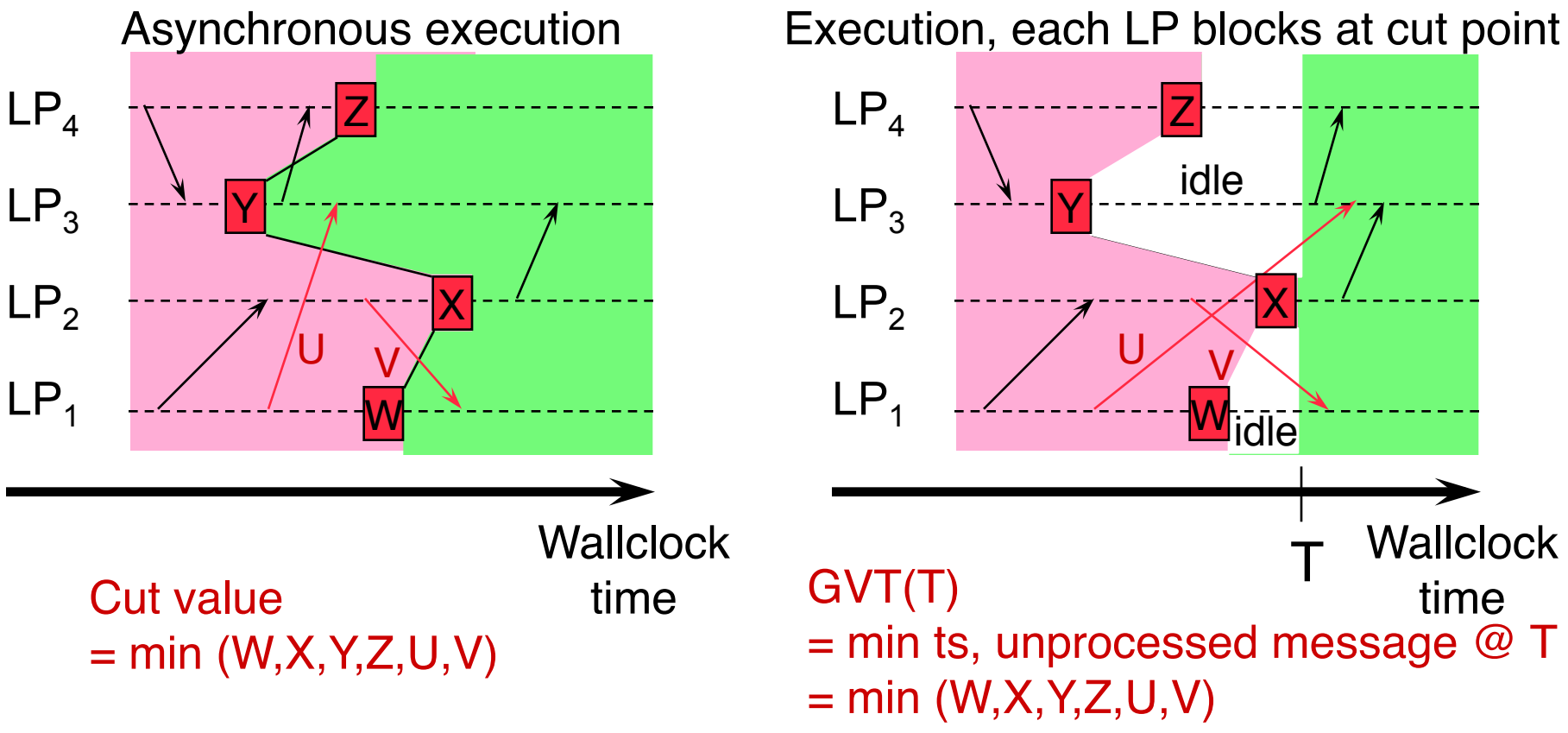
(Cut value ≤ X, Y, U)

Any message crossing cut from future to past must have a time stamp > the cut value, so they can be ignored when computing the cut value

Message generated by an LP after its cut point must have time stamp greater than the minimum of

- The LP's local minimum at its cut point
- The time stamp of messages received after the cut point

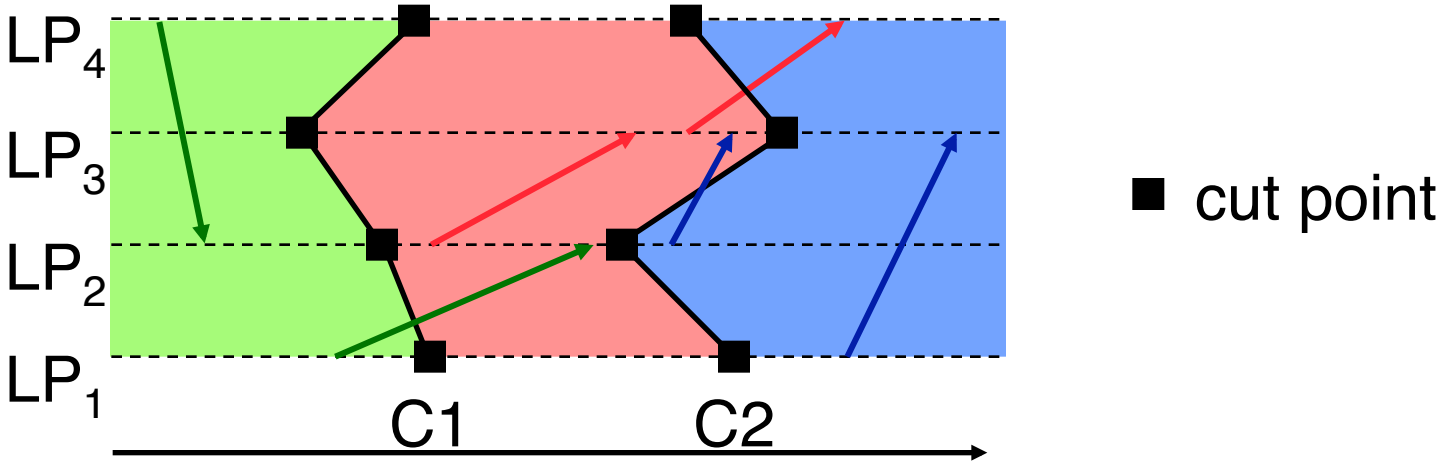
Observation 2



- Cut value equal to GVT(T) using synchronous GVT algorithm
- Events generated after cut have time stamp > cut value
- **Cut value can be used as a GVT value**

Mattern's GVT Algorithm

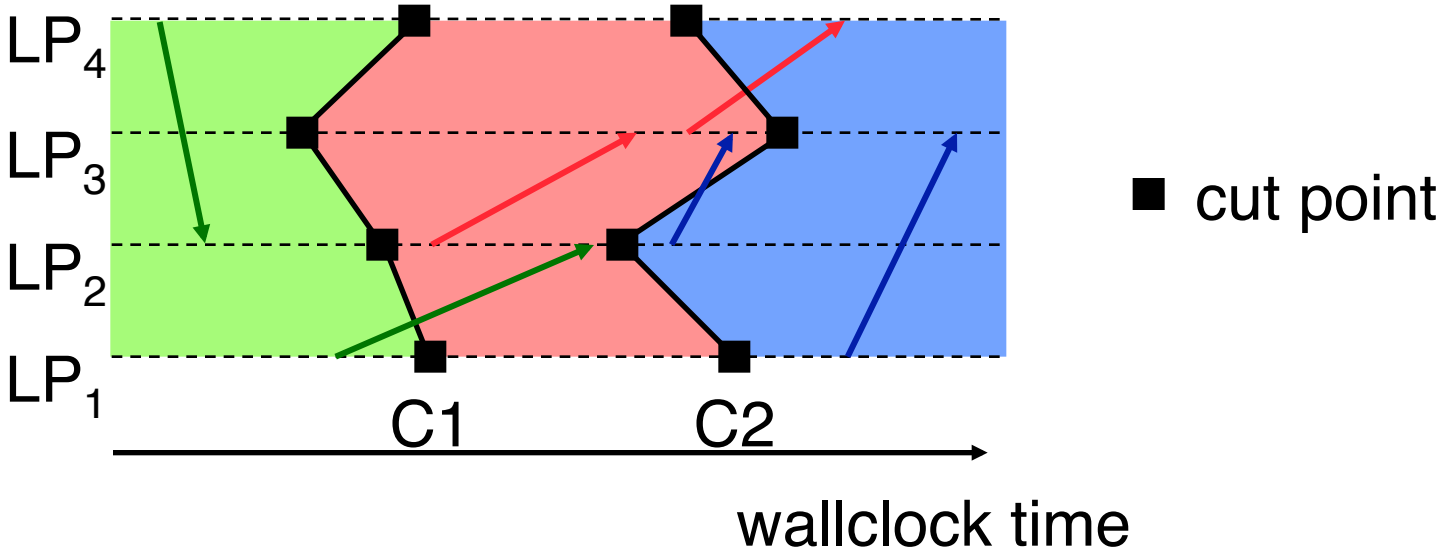
Challenge: accounting for cut messages



Approach

- Construct two cuts C1, C2, approximate cut value along C2
  - Organize processes in ring, pass token around ring
- Ensure no message that crosses C1 also cross C2
  - Color LPs, change LP color at each cut point
  - Color each message to that of LP sending message (message tag)
  - Maintain send/receive message counters
- GVT = min (local min along C2, time stamp of red messages)

Algorithm Overview



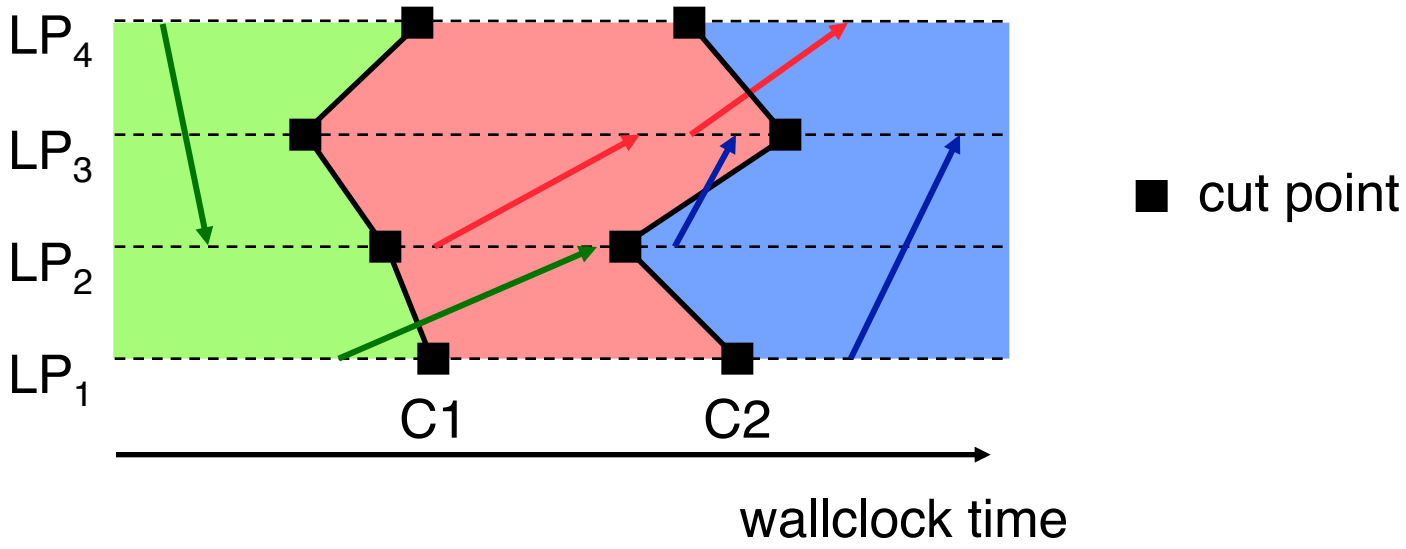
The first cut

- Changes color of each process (green to red)
- Determine number of green messages sent to each process

The second cut:

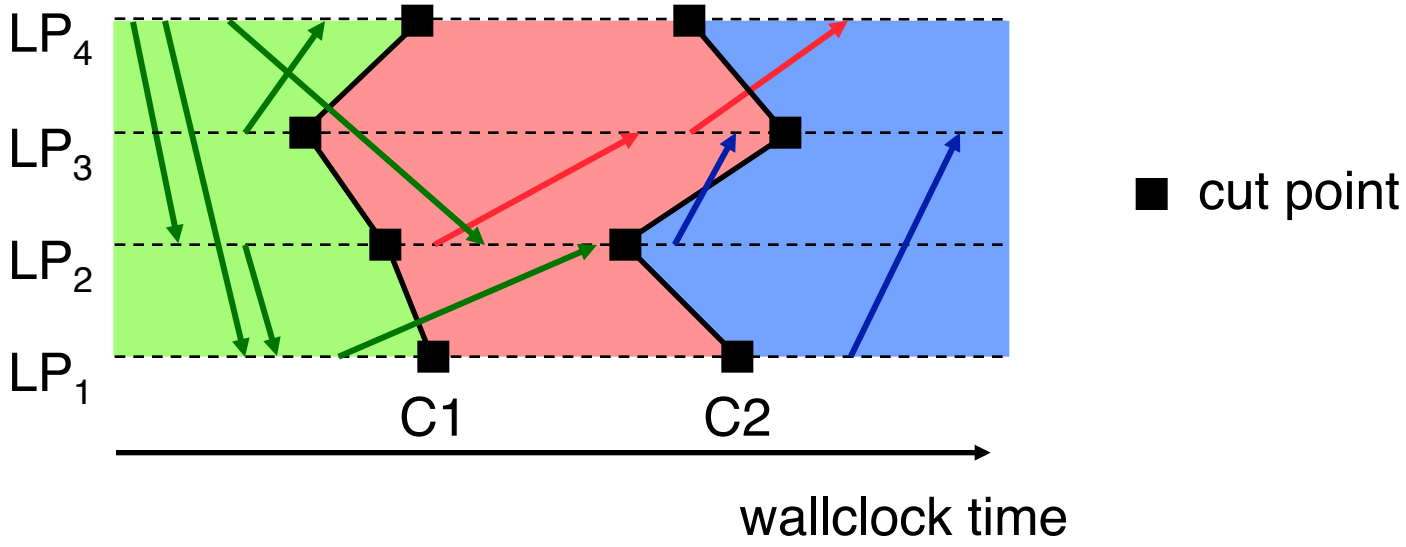
- Each process makes sure all green messages sent to it have been received before laying down a cut point
- Compute GVT: min among local mins, red messages

How an LP knows it has received all its green messages



- LP_i maintains vector $V_i[1:N]$, where N = #LPs
  - $V_i[i]$ = number of green messages received by LP_i
  - $V_i[r]$ = number of green messages sent by LP_i to LP_r
- C2: LP_i cannot pass token until
  - $V_i[i] = \sum V_s[i]$ (summed over all $s \neq i$)
- C1: Token includes vector to accumulate send counters

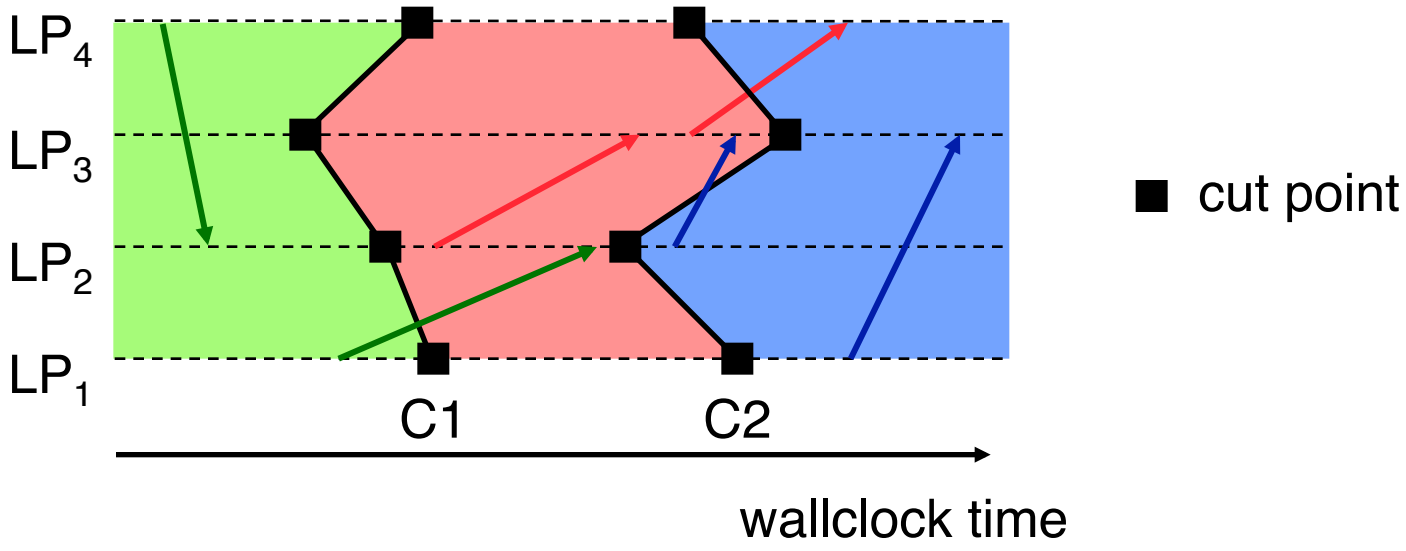
Example: Vector Counters



Vector counters for green messages (at C2):

V_1:	V_2:	V_3:	V_4:
$V_1[4]=0$	$V_2[4]=0$	$V_3[4]=1$	$V_4[4]=1$
$V_1[3]=0$	$V_2[3]=0$	$V_3[3]=0$	$V_4[3]=0$
$V_1[2]=1$	$V_2[2]=3$	$V_3[2]=0$	$V_4[2]=2$
$V_1[1]=2$	$V_2[1]=1$	$V_3[1]=0$	$V_4[1]=1$

GVT Algorithm



- Local Variables (in each logical process LP_i):
  - T_{red} = min time stamp among red messages sent by LP (even non-cut red messages!)
  - $V_i[1:N]$ = message send / receive counters
- Token: CMsg
  - $CMsg\_T_{min}$ = accumulator, smallest local minimum so far
  - $CMsg\_T_{red}$ = accumulator, smallest red message time stamp so far
  - CMsg_Count[1:N] = # messages sent to each LP

Sketch of Algorithm

- Message send by green logical process from LP_i to LP_j

  $V_i[j] = V_i[j] + 1$

- LP_i receives a green message

  $V_i[i] = V_i[i] + 1$

- Control message, first cut:

  Change color of process to red

  $CMsg\_Count = CMsg\_Count + V_i$

  Forward control message to next process in ring

- Message send with time stamp ts by a red LP

  $T_{red} = \min (T_{red}, ts)$

- Control message, Second cut:

  wait until ($V_i[i] = CMsg\_Count[i]$)

  $CMsg\_T_{min} = \min (CMsg\_T_{min}, T_{min})$

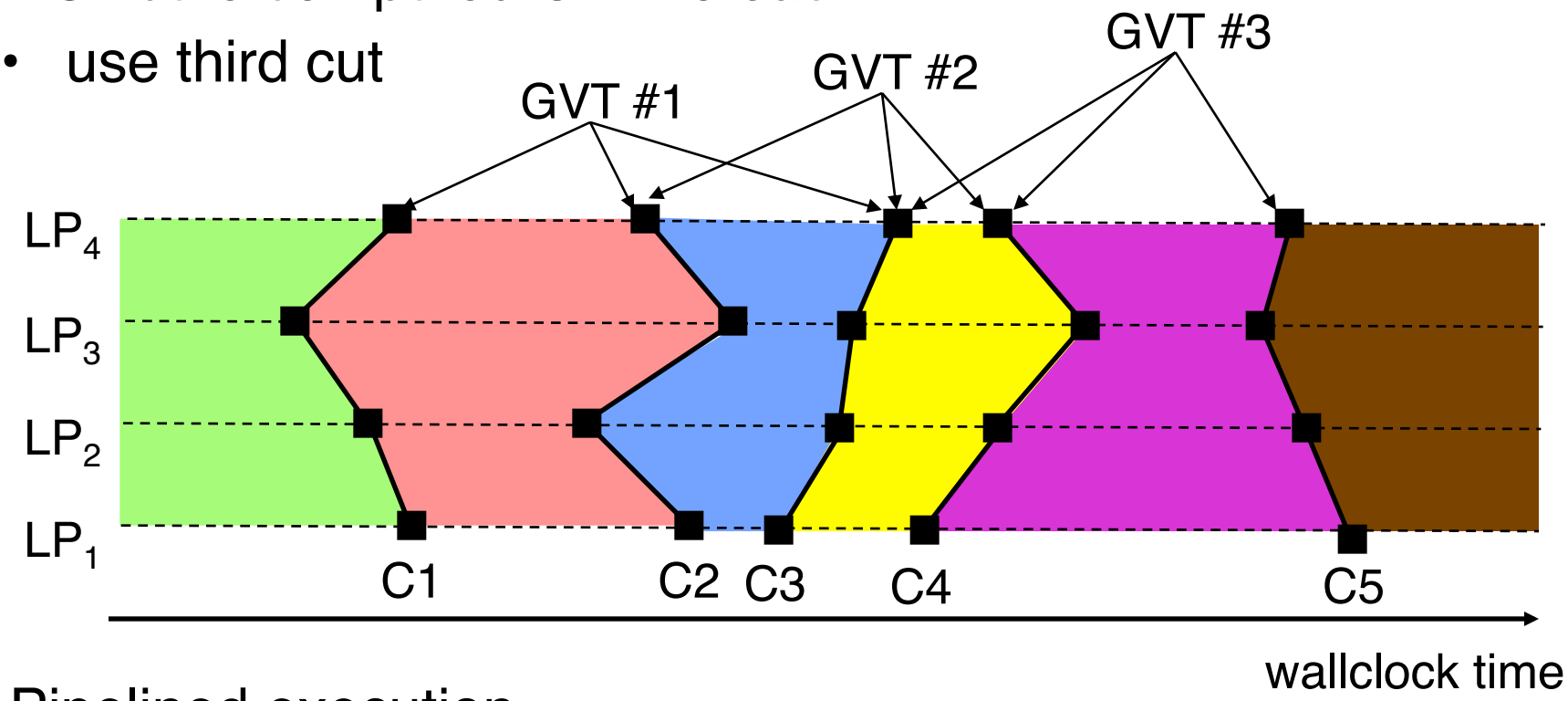
  $CMsg\_T_{red} = \min (CMsg\_T_{red}, T_{red})$

  forward token to next process in ring

Distributing GVT Values & Pipelining

Distribute computed GVT to each LP

- use third cut



Pipelined execution

- Overlap successive GVT computations: first GVT uses C1, C2, C3, second uses C2, C3, C4, etc.
- Each cut computes a new GVT value
- Continuously circulate GVT token

Fossil Collection

- Batch fossil collection
  - After GVT computation, scan through list of LPs mapped to processor to reclaim memory and commit I/O operations
  - May be time consuming if many LPs
- On-the-fly Fossil Collection
  - After processing event, place memory into “free memory” list
  - Before allocating memory, check that time stamp is less than GVT before reusing memory

Summary

- Consistent cuts
- Cut value can be used as an estimate of GVT
  - Local minimum at each LP
  - Cut messages
- Construct second consistent cut
  - Coloring LPs, messages
  - Vector counter to determine when an LP has received all relevant cut messages
- Pipeline GVT computation, continuously circulating token
- Numerous variations
  - Could implement cuts with other communication topologies, e.g., butterfly
  - Other ways to deal with transient messages, e.g., global count and abort/retry mechanism for second cut, etc.